FlowMatrix: GPU-Assisted Information-Flow Analysis through Matrix-Based Representation

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Dynamic Information Flow Tracking (DIFT)

- **DIFT** (aka **Dynamic Taint Analysis**): An important program analysis technique in security
- Track information flows in a program: Taint state transforms between sources and sinks of interest
- Security applications: Vulnerability analysis, Configuration diagnosis, etc.

```
fp = fonen(path, "r");
while(fgets(msg, maxLength-1, fp)) {
    strcat(msg, "n");
    send(socket, msg, strlen(msg), 0);
}

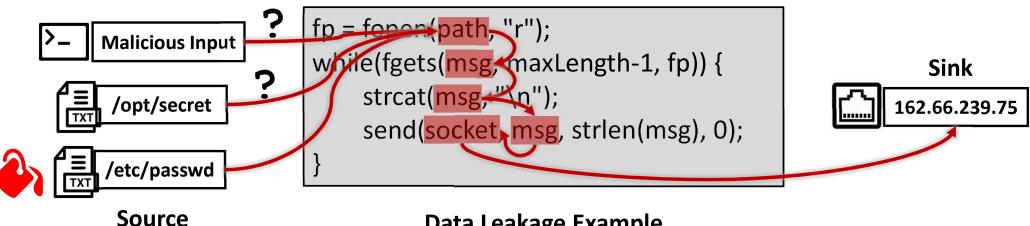
    /etc/passwd
}
```

Data Leakage Example

Source

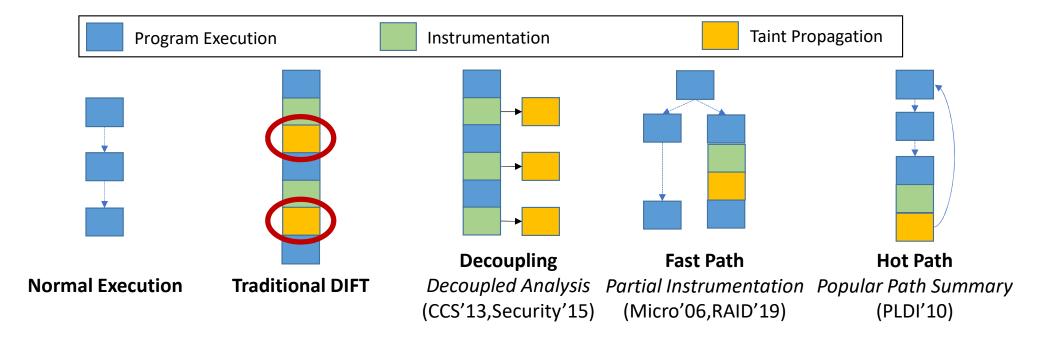
Dynamic Information Flow Tracking (DIFT)

- Challenge: Users often need to check multiple information flows
 - Calls for efficient **DIFT Query**: Rapidly DIFT with different given sources and sinks
- DIFT is **expensive**: 4~8 times performance overhead
 - One way to support DIFT query: Heavy computing support (OSDI'16)
 - Another way is to speed up DIFT itself...



Problem of DIFT

Existing work of accelerating DIFT



Lack of speeding up propagation operation itself.

Complexity of DIFT Operation Rules

- Taint propagation logics in current DIFT mechanisms are
 - Implemented in high-level programming languages with if and loops
 - Unnecessarily complex
 - Challenging to be computationally speeded-up

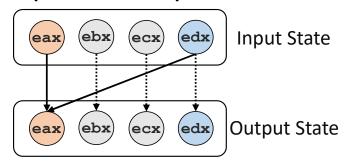
```
void taint_parallel_compute(shad, void r2r_binary_opl(dst, src, ...) int gen_taintcheck_insn(...) {
    dest, opcode, ...)
                                                                        switch(opc) {
{
                                     thread_ctx -> vcpu.gpr[dst] |=
                                                                         case INDEX_op_or_i32:
                                        thread_ctx -> vcpu.gpr[src];
                                                                          /* t0 = arg1 || arg2 */
 if (opcode == llvm::Instruction::
                                                                          tcg_gen_or_i32(t0,arg1,arg2);
    0r) {
                                                                          /* t2 = (t0 != 0)
  cb_mask_out.cb_mask =
                                    void ins_inspect(INS ins) {
   (cb_mask_1.zero_mask &
                                                                          tcg_gen_movi_i32(t_zero,0);
    cb_mask_2.cb_mask) [
                                     switch (ins_indx) {
                                                                          tcg_gen_setcond_i32(TCG_COND_NE
   (cb_mask_2.zero_mask &
                                     case XED_ICLASS_OR:
                                                                            .t2.t zero.t0):
                                                                          /* result = ~t2
    cb_mask_1.cb_mask);
                                      INS_InsertCall(
                                       r2r_binary_opl,
                                                                          tcg_gen_neg_i32(result,t2);
                                       REG32_INDX (reg_dst).
 write_cb_masks(shad, dest,
                                                                          break;
    cb_mask_out, ...);
                                       REG32_INDX(reg_src), ...);
}
```

(a) Panda (b) Libdft (c) Decaf

Different implementations of taint propagation rule of or instruction in Panda, libdft and DECAF.

Insights

- DIFT propagation logic is data dependency (TaintInduce NDSS'19)
- Example: DIFT operations for x86 instructions



```
egin{aligned} eax_{out} &= 1*eax_{in} + 1*edx_{in} \ ebx_{out} &= 1*ebx_{in} \ ecx_{out} &= 1*ecx_{in} \ edx_{out} &= 1*edx_{in} \end{aligned}
```

DIFT operations for instruction

OR eax, edx

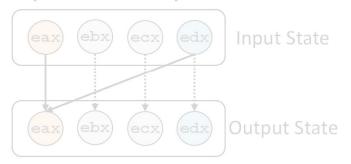
Dependencies in Boolean space

```
\begin{array}{l} eax_{out} = 1*eax_{in} + 0*ebx_{in} + 0*ecx_{in} + 1*edx_{in} \\ ebx_{out} = 0*eax_{in} + 1*ebx_{in} + 0*ecx_{in} + 0*edx_{in} \\ ecx_{out} = 0*eax_{in} + 0*ebx_{in} + 1*ecx_{in} + 0*edx_{in} \\ edx_{out} = 0*eax_{in} + 0*ebx_{in} + 0*ecx_{in} + 1*edx_{in} \end{array}
```

Dependencies in a verbose form

Insights

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- Example: DIFT operations for x86 instructions



```
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```

DIFT operations for instruction

OR eax, edx

Dependencies in Boolean space

```
eax_{out} = 1 * eax_{in} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in}
ebx_{out} = 0 * eax_{in} + 1 * ebx_{in} + 0 * ecx_{in} + 0 * eax_{in}
ecx_{out} = 0 * eax_{in} + 0 * ebx_{in} + 1 * ecx_{in} + 0 * edx_{in}
edx_{out} = 0 * eax_{in} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in}
```

Dependencies in a verbose form

Insights

We identify the <u>linearity</u> in DIFT:

 The DIFT operation between input states and output states is a linear relationship.

DIFT operations for OR eax,

A system of linear equations:

$$f: S_{in} \to S_{out}$$

$$\begin{aligned} eax_{out} &= 1 * eax_{in} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in} \\ ebx_{out} &= 0 * eax_{in} + 1 * ebx_{in} + 0 * ecx_{in} + 0 * edx_{in} \\ ecx_{out} &= 0 * eax_{in} + 0 * ebx_{in} + 1 * ecx_{in} + 0 * edx_{in} \\ edx_{out} &= 0 * eax_{in} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in} \end{aligned}$$

Dependencies in a verbose form

DIFT Operations as Matrix Transformations

FlowMatrix: a new matrix-based representation of DIFT propagation rule

```
eax_{out} = 1 * eax_{in} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in}
ebx_{out} = 0 * eax_{in} + 1 * ebx_{in} + 0 * ecx_{in} + 0 * edx_{in}
ecx_{out} = 0 * eax_{in} + 0 * ebx_{in} + 1 * ecx_{in} + 0 * edx_{in}
edx_{out} = 0 * eax_{in} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in}
edx_{out} = 0 * eax_{in} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in}
edx_{out} = 0 * eax_{out} + 0 * ebx_{in} + 0 * ecx_{in} + 1 * edx_{in}
ebx_{in} = ecx_{in}
ecx_{in} = ecx_{in}
ecx_{in} = ecx_{in}
ecx_{in} = ecx_{in}
ecx_{in} = ecx_{in}
```

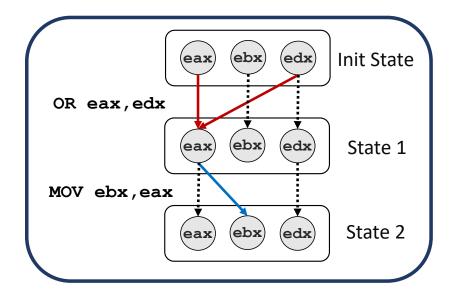
The coefficient matrix, the dependencies between S_{in} and S_{out}

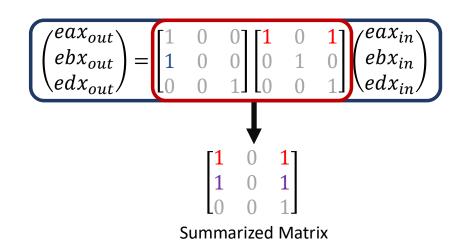
Propagation Summary as Matrix Multiplication

- Example: DIFT propagation of two x86 instructions
- Summarizing two DIFT propagation rules is to multiply two FlowMatrices:

$$M_{sum} = M_2 \times M_1$$

• FlowMatrix operations: matrix-matrix multiplication, etc.

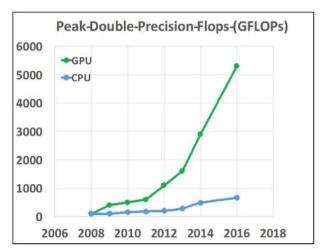


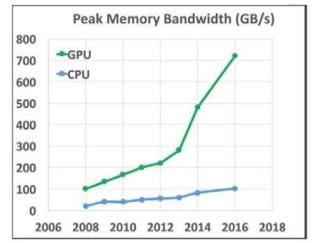


GPU-assisted DIFT Operations

- GPUs are suitable for highly parallel applications such as matrix and vector computations.
- FlowMatrix operations are accelerated by GPUs!

Speed of calculation (FLOPS) and data movement (GB/s) - #EmeringTech #MegaTrend



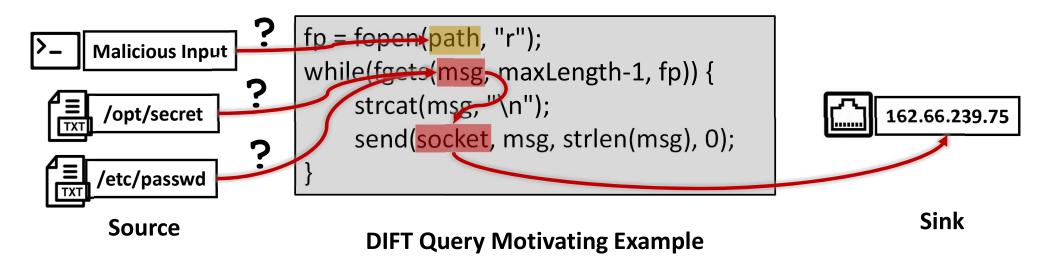


Source: HPC 2016.

source europa.eu via @mikequindazzi

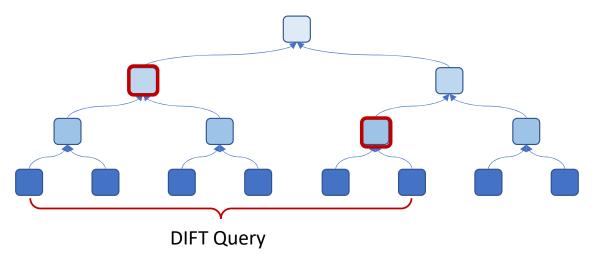
GPU-assisted FlowMatrix-based DIFT Query

- How can GPUs and FlowMatrix support efficient DIFT queries?
 - Answer a query by propagating each instruction sequentially?
 Query too slow
 - Prepare queries by pre-computing every possible query? Too much to prepare
- Goal: Reasonable pre-processing cost and rapid query response



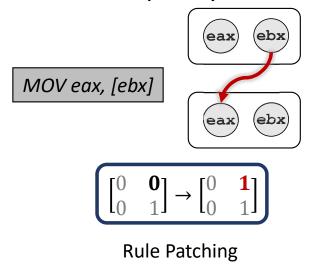
Trace-based Repeated DIFT Query

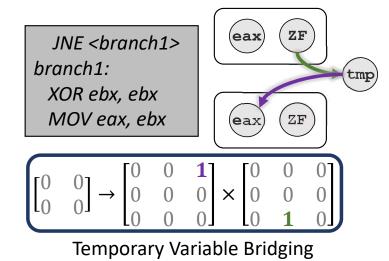
- Offline DIFT query on instruction execution traces
- (Segment-tree-like) Query Tree
 - Leaf nodes: FlowMatrix for a single instruction
 - Non-leaf nodes: Summarized FlowMatrices of two child nodes
- Pre-processing (Tree Construction): Linear time complexity
- Query: Logarithmic time complexity



Under/Over-tracking in DIFT queries

- Improper tracking policy may lead under/over-tracking
 - E.g., Common under-tracking cases: dependencies between pointers and values, between condition and in-branch variables
- How to mutate tracking policy with FlowMatrix?
 - Directly patch DIFT rule matrix
 - Add a temporary variable to bridge information flows



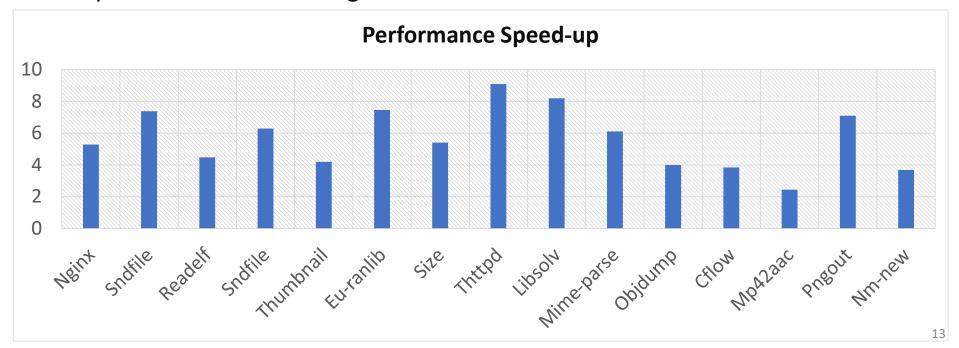


Evaluation

- Evaluation Aspects
 - Performance
 - How much improvement is achieved by GPU assistance?
 - How fast is FlowMatrix-based DIFT query?
 - Throughput
 - What is the throughput of FlowMatrix-based DIFT queries?
 - Comparison
 - How does FlowMatrix-based DIFT query compare with existing taint tools and DIFT query systems?
- Date Set
 - 15 CVEs and 7 common applications

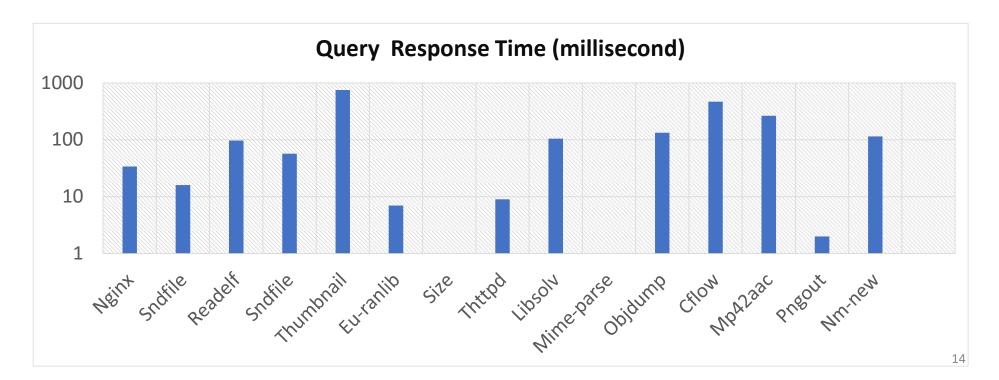
Evaluation - Performance

- Question: How much improvement achieved by GPU assistance?
- Answer:
 - Our prototype outperforms CPU-based DIFT tool over **5** times in performance on average.



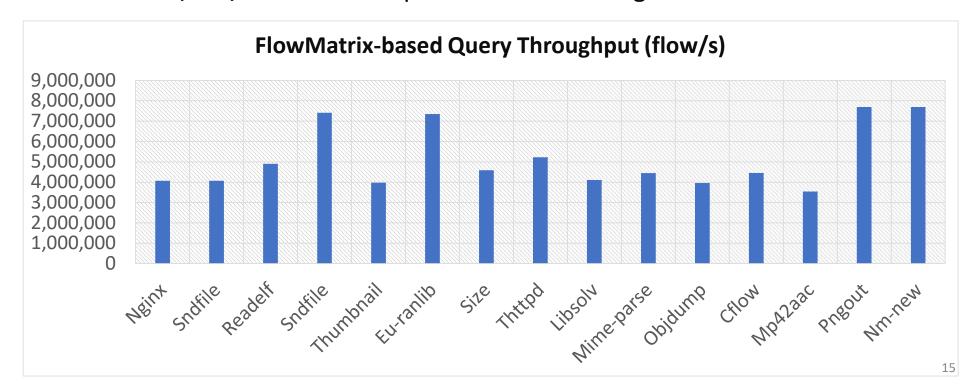
Evaluation - Performance

- Question: How fast is FlowMatrix-based DIFT query?
- Answer:
 - Most DIFT query requests can be answered in less than 0.5 sec.



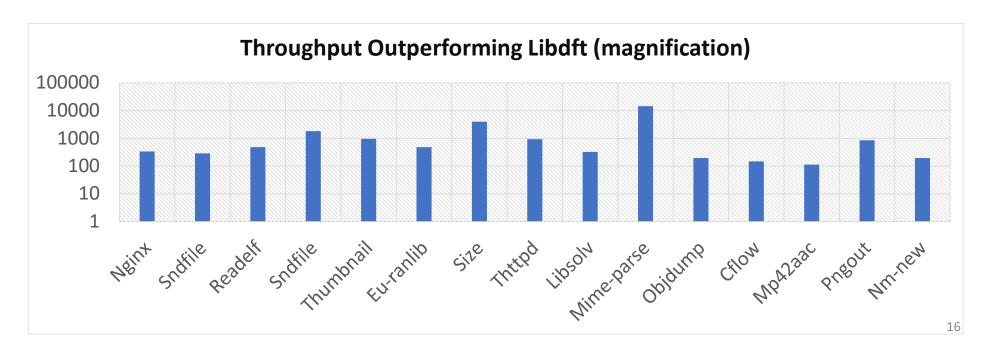
Evaluation - Throughput

- Question: What is the throughput of the DIFT query operations?
- Answer:
 - Over **5,000,000** dataflows per second on average



Evaluation - Comparison

- Question: Is FlowMatrix comparable to existing taint engines and DIFT query systems?
- Answer:
 - Three orders of magnitude larger than LibDFT
 - Comparable with JetStream (achieved by 128 CPU cores)



Summary

FlowMatrix: a Matrix-based DIFT Representation

- We recognize linearity of dynamic information flow operations
- We propose a matrix-based representation for DIFT operations

GPU-assisted DIFT

FlowMatrix enables GPU as co-processors for efficient DIFT operations

DIFT Query

We design an efficient DIFT query with high throughput

Thanks!

Q&A

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Code Available at https://github.com/mimicji/FlowMatrix